## TOWARDS LATTICE QCD+QED SIMULATIONS ON GPUS


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QCD ON THE LATTICE

## NON-ABELIAN GAUGE THEORIES

Yang-Mills Lagrangian (4D, with P and T invariance):

$$
\begin{aligned}
\mathcal{L} & =-\frac{1}{4} F_{\mu \nu}^{a} F^{\mu \nu, a}+\bar{\psi}(i \not \emptyset-m) \psi \\
\not D & =\gamma^{\mu} D_{\mu} \\
F_{\mu \nu}^{a} & =\partial_{\mu} A_{\nu}^{a}(x)-\partial_{\nu} A_{\mu}^{a}+g f^{a b c} A_{\mu}^{b}(x) A_{\nu}^{c}(x) \\
D_{\mu} & =\partial_{\mu}+i g A_{\mu}^{a}(x) T^{a},
\end{aligned}
$$

with action $\mathcal{S}=\int d^{4} x \mathcal{L}$. In path integral formalism, we have the Minkowski expectation value

$$
\begin{aligned}
\langle\mathcal{O}[\psi, \bar{\psi}, A]\rangle & =\lim _{t \rightarrow \infty(1-i \epsilon)} \frac{1}{\bar{Z}} \int \mathcal{D} \psi \mathcal{D} \bar{\psi} \mathcal{D A O}[\psi, \bar{\psi}, A] e^{i S[\psi, \bar{\psi}, A]} \\
Z & =\int \mathcal{D} \psi \mathcal{D} \bar{\psi} \mathcal{D A} e^{i S[\psi, \bar{\psi}, A]} .
\end{aligned}
$$

## DISCRETIZATION

4D lattice

$$
\Lambda=\left\{n=\left(n_{0}, n_{1}, n_{2}, n_{3}\right) \mid n_{i} \in\left\{0,1, \ldots, L_{i}-1\right\}, i \in\{0,1,2,3\}\right\}
$$

$$
\begin{aligned}
& A_{\mu}^{b}(x)=A_{\mu}^{b}(a n) \longrightarrow a A_{\mu}^{b}(n), \\
& \psi(x)=\psi(a n) \longrightarrow a^{3 / 2} \psi(n) \\
& U(x, x+\epsilon \hat{\mu})=U(a n, a n+a \hat{\mu}) \longrightarrow U_{\mu}(n) \text {. } \\
& \text { - lattice extents } L_{i} \in \mathbb{N} \\
& \text { - lattice constant } a \\
& \text { - compensator field (link variable) } \\
& U(x, y)\left(\frac{\partial U}{\partial x^{\mu}}=i g A_{\mu}^{a}(x) T^{a}\right)
\end{aligned}
$$

## DISCRETIZATION

The full discretized Yang-Mills action is

$$
\begin{aligned}
\mathcal{S} & =\mathcal{S}_{G}+\mathcal{S}_{F}, \\
\mathcal{S}_{G} & =\frac{1}{g^{2}} \sum_{n \in \Lambda} \sum_{\mu, \nu} \operatorname{Re} \operatorname{tr}\left[\delta_{\mu \nu} \cdot i d-\hat{U}_{\mu \nu}(n)\right], \\
\mathcal{S}_{F} & =a^{4} \sum_{n \in \Lambda} \bar{\psi}(n) \underbrace{\left[\sum_{\mu} \frac{\gamma^{\mu}}{2}\left(D_{+\mu}+D_{-\mu}\right)+m\right]}_{D} \psi(n)
\end{aligned}
$$




- $D_{ \pm \mu}$ are forward- and backward covariant derivatives
- $D$ is the Dirac operator
- Large sparse matrix, $\operatorname{dim} d=12 \mathrm{~V}$, (eg. $64^{4}$ lattice: $d \approx 10^{8}$, naively 500 petabytes)


## WICK ROTATION

Rotating to the Minkowski action $e^{i S} \xrightarrow{\text { WR }} e^{-\mathcal{S}_{E}}$ we obtain the Euclidean expectation value

$$
\begin{aligned}
\langle\mathcal{O}[\psi, \bar{\psi}, A]\rangle & =\lim _{\tau \rightarrow \infty} \frac{1}{\bar{Z}} \int \mathcal{D} \psi \mathcal{D} \bar{\psi} \mathcal{D A O}[\psi, \bar{\psi}, A] e^{-S[\psi, \bar{\psi}, A]} \\
Z & =\int \mathcal{D A} \operatorname{det}(D) e^{-S_{G}[A]} .
\end{aligned}
$$

Interpret the exponential as probability density, $P(U)=Z^{-1} e^{-S}$,

$$
\begin{aligned}
\langle\mathcal{O}[U]\rangle & =\lim _{\tau \rightarrow \infty} \frac{1}{Z} \int \mathcal{D} U e^{-S[U]} \mathcal{O}[U] \longrightarrow \sum_{U} P(U) \mathcal{O}[U] \\
\langle\mathcal{O}[U]\rangle & \approx \frac{1}{N} \sum_{i=1}^{N} \mathcal{O}\left[U_{i}\right] .
\end{aligned}
$$

## LATTICE QFT ON THE CPU



- large sparse linear operator (complex dimension $O\left(10^{8}\right)$ )
- Markov chain alters using Hybrid Monte Carlo [1]

■ operator needs to be inverted repeatedly (solve $D \psi=\eta$ )

- most used compute intensive kernel: sparse matrix-vector multiplication (SpMV), D $\psi$
- sparse linear algebra is memory bound (speed limited by memory bandwidth, not compute power)
- using supercomputers ( $O(100$ ) nodes with $O(100)$ cores) $\rightarrow$ high parallelizability
- large problem is decomposed into many smaller problems


## OPENQXD

## 

■ based on openQCD 1.6 written by Martin Lüscher and Stefan Schäfer [2]
■ implementing dynamical QCD+QED simulations

- addition of $C^{*}$ boundary conditions

■ available via https://gitlab.com/rcstar/openQxD
■ MPI 1.2, ISO C90
■ currently no threading, GPU support

DATA TYPES

## FloAting Point formats

| Floating-point formats |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| name | s | e | m | machine $\epsilon\left(2^{-m}\right)$ |
| binary64 [3] | 1 | 11 | 52 | $2.2 \times 10^{-16}$ |
| binary32 | 1 | 8 | 23 | $1.2 \times 10^{-7}$ |
| binary16 | 1 | 5 | 10 | $9.8 \times 10^{-4}$ |
| bfloat16[4] | 1 | 8 | 7 | $7.8 \times 10^{-3}$ |
| tensorfloat32[5] | 1 | 8 | 10 | $9.8 \times 10^{-4}$ |
| posit32 [6] | 1 | es=2 |  |  |
| posit16 | 1 | es=1 |  | $7.5 \times 10^{-9}$ |
| posit8 | 1 | es=0 |  | $3.1 \times 10^{-4}$ |

Table: Commonly used floating-point formats, where $s$ is the number of sign bits, $e$ the number of exponent bits and $m$ the number of mantissa bits.


## FLOAtS [3]



■ $f=(-1)^{S} \cdot M \cdot 2^{E}$

- two representations of o
- $-\infty$ and $+\infty$
- lots of NaNs
- may over- or underflow
- complicated comparison operations
- existence of subnormals

■ many cases $\Longrightarrow$ large on-chip FPU $\Longrightarrow$ small operations per watt count Examples:

○ ๑...○ ๑... $0=+0$
1 ○...0 $0 . . .0=-0$

- 1...1 $0 . . .0=+\infty$

1 1...1 $0 . . .0=-\infty$

## Posits [6]

sign


■ $p=(-1)^{S} \cdot\left(2^{2^{e s}}\right)^{k} \cdot M \cdot 2^{E}$
■ no NaNs, but calculation is interrupted
■ only one representation of o

- every number has a reciprocal
- reciprocal of $o$ is $\pm \infty$

■ may not over- or underflow
■ simple comparison (bits identical $\equiv$ numbers equal)
■ Posit processing unit takes less circuitry than an FPU $\Longrightarrow$ higher operations per watt count

## FLOATS VS. POSITS



CONJUGATE GRADIENT

## CONJUGATE GRADIENT - CONVERGENCE ANALYSIS



Conclusions:

- binary16 or bfloat16 are sufficient
■ use binary64 in reduction variables (norms)

$$
\|\vec{x}\|=\sqrt{\sum_{i} x_{i}^{2}}
$$

■ implement a general mixed precision solver

■ solver-kernel consists of norms, scalar products, applications of D, axpys $\Longrightarrow$ memory bound operations
■ Residue: $\rho_{i}=\eta-D \psi_{i}$

- binary16: 5 exponent, 10 mantissa bits

■ bfloat16: 8 exponent, 7 mantissa bits

SAP-GCR ALGORITHM

## SAP-GCR ALGORITHM

- Decompose the lattice into many blocks
- Only nearest-neighbour interaction $\Longrightarrow$ blocks of same color are independent
- solve gray blocks
- update boundaries
- solve white blocks
- update boundaries
$\Longrightarrow$ one Schwarz-cycle (alternate between black and white blocks)
■ Preconditioning phase:
- $n_{c y}$ Schwarz-cycles
- $n_{m r}$ MR-steps on each blocked problem

| $\Omega_{1}$ | $\Omega_{2}$ | $\Omega_{3}$ | $\Omega_{4}$ |
| :--- | :--- | :--- | :--- |
| $\Omega_{5}$ | $\Omega_{6}$ | $\Omega_{7}$ | $\Omega_{8}$ |
| $\Omega_{9}$ | $\Omega_{10}$ | $\Omega_{11}$ | $\Omega_{12}$ |
| $\Omega_{13}$ | $\Omega_{14}$ | $\Omega_{15}$ | $\Omega_{16}$ |

Figure: A two dimensional example of a decomposition of a lattice $\Omega=\bigcup_{i=1}^{n} \Omega_{i}$ into $n=16$ domains $\Omega_{i}[7]$.

## MEASUREMENTS



Figure: Time measurements for the SAP_GCR kernel on different matrices and configurations. The measurements were conducted on an AMD EPYC 7742 CPU @ 2.25 GHz with 512 GB memory and an NVIDIA A100 (via SXM4) GPU with 40 GB memory.

## SAP+GCR - CONCLUSIONS

- Heavy and non-intuitive run-time dependence on input parameters $\left(n_{c y}, n_{m r}\right)$
■ Existence of exceptional configurations with extremely long run-times, non-convergence
■ Adaptive variant (tries to find the optimal config every GCR-iteration anew, avoiding exceptional configs, suitable long running simulations, where D vastly changes its condition)
- upper bound: $n_{c y}=20$ and $n_{m r}=20$
- lower bound: $n_{c y}=1$ and $n_{m r}=4$
- after every Schwarz cycle, exit if residual satisfies

$$
\left\|\rho_{i}\right\| \geq\left\|\rho_{i-1}\right\|
$$

- after every MR-step, exit if blocked residual satisfies

$$
\left\|\rho_{i}\right\| \geq \alpha\left\|\rho_{i-1}\right\| \text { where } \alpha \in\{0.7,0.9\}
$$

## SAP-GCR SOLVER BASED ON INEXACT DEFLATION

## DFL SAP GCR

- Dirac operator has a property called local coherence [8]
$\rightarrow$ mode modes look "the same" locally
$\rightarrow$ few low modes are enough to build large deflation subspaces by block projection
■ Second stage of preconditioning
- $D$ restricted to the deflation subspace: little Dirac operator A
- In every outer GCR step we solve the little
 equation


## SAP-GCR SOLVER BASED ON INEXACT DEFLATION



Figure: Time measurements for the DFL_SAP_GCR kernel on different matrices and configurations. The measurements were conducted on an Intel(R) 6130 @ 2.10 GHz with 1.5 TB memory and an NVIDIA V100 (via PCIe) GPU with 16 GB memory.

SUMMARY AND CONCLUSIONS

## SUMMARY AND CONCLUSIONS

■ Convergence analysis of CG using different data types

- Mixed precision solver desirable, specially employing 16-bit data types (i.e. bfloat16 on NVIDIA A100)
- Mild benefits from posits as opposed to floats
- These insights apply to the other solvers as well
- SAP-GCR solver
- Sensitive to choice of input parameters
- Adaptive variant might be beneficial

■ Deflated SAP-GCR solver

- Non-trivial interplay between deflation and SAP

OUTLOOK AND CURRENT STATE

## OUTLOOK AND CURRENT STATE

Motivated by the above analysis
■ Successfully coupled openQxD to QUDA [9]

- Gauge field transfer
- Multi-GPU support
- Still missing steps to tackle:
- Fermion field transfer
- C* boundaries
- QCD+QED


## Thank you for listening!

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